

Tutorial

The Perturbation-Duality Scheme in Optimization

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February 6, 2025



Introduction

Duality is widely used in optimization

- ▶ linear programming
(Lagrangian duality, including optimal transport, etc.)
- ▶ convex programming
(Lagrangian duality in mathematical programming,
minimal cost flow on a graph etc.)
- ▶ conic programming, semidefinite programming, etc.

The **perturbation-duality scheme** (PDS)

- ▶ Introduced in [Rockafellar, 1974]
- ▶ Goal: systematically produce dual optimization problems from a given optimization problem by perturbation followed by conjugate duality

Outline

PDS for linear programming (LP)

PDS for convex programming

PDS for pure integer linear programming (PILP)

Conclusion

Outline of the presentation

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Fundamental LP duality theorem

LP weak duality through PDS

LP strong duality through PDS

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Primal and dual problems in standard LP

$$\begin{array}{ccc} \text{dual problem} & & \text{primal problem} \\ \hline \sup & \langle b \mid \lambda \rangle & \inf \quad k^T x \\ \lambda \in \mathbb{R}^m & & x \in \mathbb{R}^n \\ \lambda^T A \leq k & \underbrace{\leq}_{\text{weak duality}} & Ax = b \\ & & x \geq 0 \end{array}$$

Strong duality for LP

Adapted from [Conforti, Cornuéjols, and Zambelli, 2014, Theorem 3.7, Proposition 3.9]

Theorem

Given a matrix $A \in \mathbb{R}^{m \times n}$ and vectors $k \in \mathbb{R}^n, b \in \mathbb{R}^m$,
if $\{x \in \mathbb{R}^n \mid Ax = b, x \geq 0\} \neq \emptyset$ or $\{\lambda \in \mathbb{R}^m \mid \lambda^T A \leq k\} \neq \emptyset$,
(that is, if the primal or the dual problem is feasible)

then we have

$$\begin{array}{ccc} \sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T A \leq k}} \langle b \mid \lambda \rangle & \underbrace{=} & \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} k^T x \end{array}$$

strong duality

Sketch of the proof

1. Introduce the **Lagrangian** $\mathcal{L}: \mathbb{R}_+^n \times \mathbb{R}^m \rightarrow \mathbb{R}$ by

$$\mathcal{L}(x, \lambda) = k^T x + \langle b - Ax \mid \lambda \rangle, \quad \forall x \in \mathbb{R}_+^n, \lambda \in \mathbb{R}^m$$

2. Use **sup-inf inversion** inequality to get weak duality

$$\underbrace{\sup_{\lambda \in \mathbb{R}^m} \inf_{x \in \mathbb{R}_+^n} \mathcal{L}(x, \lambda)}_{\text{dual problem}} \leq \underbrace{\inf_{x \in \mathbb{R}_+^n} \sup_{\lambda \in \mathbb{R}^m} \mathcal{L}(x, \lambda)}_{\text{primal problem}}$$

3. Find a **saddle-point** $(\bar{x}, \bar{\lambda}) \in \mathbb{R}_+^n \times \mathbb{R}^m$

$$\mathcal{L}(\bar{x}, \lambda) \leq \mathcal{L}(\bar{x}, \bar{\lambda}) \leq \mathcal{L}(x, \bar{\lambda}), \quad \forall x \in \mathbb{R}_+^n, \lambda \in \mathbb{R}^m$$

to prove strong duality, i.e.

$$\underbrace{\sup_{\lambda \in \mathbb{R}^m} \inf_{x \in \mathbb{R}_+^n} \mathcal{L}(x, \lambda)}_{\text{dual problem}} = \underbrace{\inf_{x \in \mathbb{R}_+^n} \sup_{\lambda \in \mathbb{R}^m} \mathcal{L}(x, \lambda)}_{\text{primal problem}}$$

Now, let us deduce the previous duality results
from a perturbation duality scheme (PDS)

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Steps of the perturbation-duality scheme

[Rockafellar, 1974]

1. **Perturb** a minimization problem with a perturbation (primal) variable belonging to a vector space, and a Rockafellian function
2. **Pair** the (primal) perturbation space with a dual space by means of a bilinear form $\langle | \rangle$
3. **Biconjugate** the perturbation function, and get
 - ▶ a dual problem
 - ▶ weak duality
4. Deduce conditions for **strong duality** by means of either global or local properties of the perturbation function

Illustration of the scheme in Linear Programming (LP)

- ▶ Constraint matrix $A \in \mathbb{R}^{m \times n}$
- ▶ Cost vector $k \in \mathbb{R}^n$
- ▶ Anchor $\bar{b} \in \mathbb{R}^m$

Initial/original minimization problem

$$\begin{aligned} \inf \quad & k^T x \\ x \in \quad & \mathbb{R}^n \\ Ax = \quad & \bar{b} \\ x \geq \quad & 0 \end{aligned}$$

Step 1. Perturbation of the initial minimization problem

- ▶ Introduce a perturbation space, \mathbb{R}^m , and embed the original problem into a family of minimization problems (more on the Rockafellian later)
- ▶ Introduce the perturbation function

$$\varphi: \mathbb{R}^m \rightarrow \overline{\mathbb{R}} = \mathbb{R} \cup \underbrace{\{-\infty\}}_{\text{unbounded}} \cup \underbrace{\{+\infty\}}_{\text{unfeasible}}$$

$$\forall b \in \mathbb{R}^m, \quad \varphi(b) = \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} k^T x$$

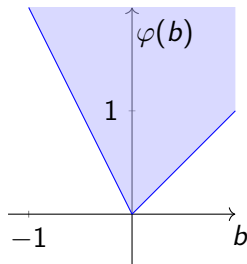
- ▶ The value of the original problem is then $\varphi(\bar{b})$

Example of perturbation function's epigraph for LP

Let $\varphi: \mathbb{R} \rightarrow \mathbb{R}$ be defined as

$$\forall b \in \mathbb{R}, \quad \varphi(b) = \inf_{\substack{x \in \mathbb{R}^2 \\ x_1 - x_2 = b \\ x \geq 0}} x_1 + 2x_2$$

Then $\varphi(b) = \max\{-2b, b\}$, $\forall b \in \mathbb{R}$



Step 2. Dual space, coupling and conjugate function

- ▶ Perturbation space: \mathbb{R}^m — dual space: \mathbb{R}^m (linear functions)
- ▶ Introduce the bilinear coupling

$$\langle | \rangle : \overbrace{\mathbb{R}^m}^{\text{perturbation space}} \times \overbrace{\mathbb{R}^m}^{\text{dual space}} \rightarrow \mathbb{R}$$

- ▶ Deduce the **conjugate function** $\varphi^*: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$ of the perturbation function

$$\forall \lambda \in \mathbb{R}^m, \quad \varphi^*(\lambda) = \sup_{b \in \mathbb{R}^m} \{ \langle b | \lambda \rangle - \varphi(b) \}$$

Conjugate function and Lagrangian

$$\begin{aligned}\varphi^*(\lambda) &= \sup_{b \in \mathbb{R}^m} \{ \langle b \mid \lambda \rangle - \varphi(b) \} \\&= \sup_{b \in \mathbb{R}^m} \left\{ \langle b \mid \lambda \rangle - \inf_{\substack{Ax=b \\ x \geq 0}} k^T x \right\} \\&= \sup_{b \in \mathbb{R}^m} \left\{ \langle b \mid \lambda \rangle + \sup_{\substack{Ax=b \\ x \geq 0}} \langle -x \mid k \rangle \right\} \\&= \sup_{x \geq 0} \left\{ \sup_{\substack{Ax=b \\ b \in \mathbb{R}^m}} \langle b \mid \lambda \rangle - k^T x \right\} \\&= \sup_{x \geq 0} \{ \langle Ax \mid \lambda \rangle - k^T x \} \\&= \langle \bar{b} \mid \lambda \rangle - \inf_{x \geq 0} \underbrace{\{ k^T x + \langle \bar{b} - Ax \mid \lambda \rangle \}}_{\text{Lagrangian } \mathcal{L}(x, \lambda)}\end{aligned}$$

Step 3. Biconjugate and weak duality

- Biconjugate function $\varphi^{**'}: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\forall b \in \mathbb{R}^m, \quad \varphi^{**'}(b) = \sup_{\lambda \in \mathbb{R}^m} \{ \langle b \mid \lambda \rangle - \varphi^*(\lambda) \}$$

- We obtain **weak duality** for all $b \in \mathbb{R}^m$

$$\underbrace{\sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T A \leq k}} \langle b \mid \lambda \rangle}_{\text{dual problem}} = \varphi^{**'}(b) \leq \varphi(b) = \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} k^T x$$

- At the anchor \bar{b}

$$\varphi^{**'}(\bar{b}) = \sup_{\lambda \in \mathbb{R}^m} \underbrace{\{ \langle \bar{b} \mid \lambda \rangle - \varphi^*(\lambda) \}}_{\inf_{x \geq 0} \mathcal{L}(x, \lambda)}$$

We have obtained the LP weak duality result
What about the LP strong duality result?

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Step 4. Conditions for strong duality

Proposition

Let $A \in \mathbb{R}^{m \times n}$ and $k \in \mathbb{R}^n$

If $-\infty < \varphi(0)$, that is, **the corresponding LP is bounded**
then, **for all $b \in \mathbb{R}^m$** ,

$$\left(\sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T A \leq k}} \langle b \mid \lambda \rangle \right) = \underbrace{\varphi^{**'}(b) = \varphi(b)}_{\text{strong duality}} \left(= \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} k^T x \right)$$

Remark

*This result is true even if $b \in \mathbb{R}^m$ is such that $\varphi(b) = +\infty$,
meaning for any unfeasible LPs*

Proof of strong duality for LP. Sketch of the proof

adapted from [Rockafellar, 1974, p.24]

- (a) We show that if the LP $\varphi(0)$ is bounded, then every feasible LP is bounded

$$-\infty < \varphi(0) \implies \varphi \text{ is proper}$$

- (b) We show that $\text{epi } \varphi$ is a closed convex set (by showing that $\text{epi } \varphi$ is a polyhedron)

$$\varphi \text{ is a closed convex function}$$

- (c) We apply Fenchel-Moreau Theorem to get strong duality

$$\varphi^{**'} = \varphi$$

Strong duality for LP. Step (a) Proper functions

Definition

Let $f : \mathbb{R}^m \rightarrow \overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty, +\infty\}$

- ▶ $\text{dom } f = \{b \in \mathbb{R}^m : f(b) < +\infty\}$
- ▶ The function f is said to be **proper** if $\text{dom } f \neq \emptyset$ and $-\infty < f(b)$, $\forall b \in \mathbb{R}^m$

Lemma

If $-\infty < \varphi(0)$ (the corresponding LP is bounded)

*then the value function φ is **proper** (all feasible LPs are bounded)*

Idea of the proof

The recession cone of $\{x \in \mathbb{R}^n : Ax = b\}$

is given by $\{r \in \mathbb{R}^n : Ax = 0, r \geq 0\}$

[Conforti, Cornuéjols, and Zambelli, 2014, Proposition 3.15]

Strong duality for LP. Step (b) Closed convex epigraph

Definition

Let $f : \mathbb{R}^m \rightarrow \overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty, +\infty\}$

The epigraph of the function f is defined by

$$\text{epi } f = \{(b, t) \in \mathbb{R}^m \times \mathbb{R} : f(b) \leq t\}$$

Proposition

Let $A \in \mathbb{R}^{m \times n}$ and $k \in \mathbb{R}^n$ define the value function $\varphi : \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$ by

$$\forall b \in \mathbb{R}^m, \quad \varphi(b) = \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} k^T x$$

Then $\text{epi } \varphi$ is a polyhedron

Proof that $\text{epi } \varphi$ is a polyhedron

$$A \in \mathbb{R}^{m \times n}, k \in \mathbb{R}^n$$

Let $b \in \mathbb{R}^m$, we assume that $-\infty < \varphi(b)$

$$\varphi(b) \leq t$$

$$\iff \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} k^T x \leq t$$

$$\iff \min_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} k^T x \leq t \quad (\text{as bounded feasible LPs are attained})$$

$$\iff \exists x \in \mathbb{R}^n \text{ s.t. } Ax = b, \ x \geq 0, \ k^T x - t \leq 0$$

$$\iff \text{epi } \varphi = \pi_{(b,t)} \left\{ (b, t, x) \in \mathbb{R}^m \times \mathbb{R} \times \mathbb{R}^n : \begin{cases} Ax = b \\ x \geq 0 \\ k^T x - t \leq 0 \end{cases} \right\}$$

Thus $\text{epi } \varphi$ is the projection of a polyhedron

So, $\text{epi } \varphi$ is a polyhedron

[Rockafellar, 1970, Theorem 19.3]

Strong duality for LP. Step (c) Fenchel-Moreau theorem

Definition

A function $f : \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$ is said to be **closed convex** [Rockafellar, 1974] if **EITHER** [f is proper **AND** $\text{epi } f$ is a closed convex set] **OR** $f \equiv +\infty$ **OR** $f \equiv -\infty$

Theorem

[Fenchel-Moreau Theorem]

A function $f : \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$ is closed convex if and only if $f^{**'} = f$

So we have **strong duality**

$$\varphi^{**'} = \underbrace{\varphi}$$

as Steps (a) and (b) imply
that φ is a closed function

About closed convex functions: the case of valley functions

Definition

Let $C \subset \mathbb{R}^n$ be a closed convex set

Let $f: \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ be a function

We say that f is a **valley function** if

$$f(u) = \begin{cases} -\infty & \text{if } u \in C \\ +\infty & \text{otherwise} \end{cases}$$

Remark

Valley functions have a closed convex epigraph

BUT *are not closed convex functions*

(except the cases $f \equiv -\infty$ or $f \equiv +\infty$)

Example when strong duality is not achieved for LP

$$-\infty = \left(\begin{array}{cc} \sup & \lambda_1 \\ \lambda \in \mathbb{R}^2 & \\ \lambda_1 + \lambda_2 \leq -1 & \\ -\lambda_1 - \lambda_2 \leq 0 & \end{array} \right)$$

$$= \varphi^{**'}((1, 0)) < \varphi((1, 0)) =$$

$$\left(\begin{array}{cc} \inf & -x_1 \\ x \in \mathbb{R}^2 & \\ x_1 - x_2 = 1 & \\ x_1 - x_2 = 0 & \\ x \geq 0 & \end{array} \right) = +\infty$$

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Summary of the perturbation-duality scheme for LP

[Rockafellar, 1974]

1. We **perturb** a minimization problem

$$\forall \mathbf{b} \in \mathbb{R}^m, \varphi(\mathbf{b}) = \inf_{\substack{\mathbf{x} \in \mathbb{R}^n \\ \mathbf{Ax} = \mathbf{b} \\ \mathbf{x} \geq 0}} \mathbf{k}^T \mathbf{x}$$

2. We pair the primal space \mathbb{R}^m and a dual space \mathbb{R}^m

$$\langle | \rangle : \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$$

3. We **biconjugate** the perturbation function φ

$$\left(\sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T \mathbf{A} \leq \mathbf{k}}} \langle \mathbf{b} | \lambda \rangle \right) = \underbrace{\varphi^{**'}(\mathbf{b}) \leq \varphi(\mathbf{b}), \forall \mathbf{b} \in \mathbb{R}^m}_{\text{weak duality is guaranteed}}$$

4. Under suitable assumptions, **strong duality** by polyhedral property of the epigraph of the perturbation function

Indicator function of a subset

For any subset $X \subset \mathcal{X}$, its indicator function ι_X is

$$\iota_X(x) = \begin{cases} 0 & \text{if } x \in X \\ +\infty & \text{if } x \notin X \end{cases}$$

Summary (bis) of the perturbation-duality scheme for LP

1. We **perturb** a minimization problem

$$\varphi(b) = \inf_{x \in \mathbb{R}^n} \mathcal{R}(x, b), \quad \forall b \in \mathbb{R}^m$$

where the **Rockafellian** $\mathcal{R}: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$ is defined by $\mathcal{R}(x, b) = k^T x + \iota_{\mathbb{R}_+^n}(x) + \iota_{\{0\}}(Ax - b)$

2. We pair the primal space \mathbb{R}^m and a dual space \mathbb{R}^m

$$\langle | \rangle : \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$$

3. We **biconjugate** the perturbation function φ

$$\left(\sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T A \leq k}} \langle b | \lambda \rangle \right) = \underbrace{\varphi^{**'}(b) \leq \varphi(b)}_{\text{weak duality is guaranteed}}, \quad \forall b \in \mathbb{R}^m$$

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- PDS with Fenchel duality

- Examples of PDS for convex programs

- Generalized perturbation duality scheme

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The Fenchel conjugacy

Definition

Two vector spaces \mathcal{U} and \mathcal{V} , paired by a bilinear form $\langle \cdot, \cdot \rangle$ (in the sense of convex analysis), give rise to the classic **Fenchel conjugacy** between $\overline{\mathbb{R}}^{\mathcal{U}}$ and $\overline{\mathbb{R}}^{\mathcal{V}}$

With any function $f: \mathcal{U} \rightarrow \overline{\mathbb{R}}$, we associate the function $f^*: \mathcal{V} \rightarrow \overline{\mathbb{R}}$ defined by

$$f^*(v) = \sup_{u \in \mathcal{U}} \left\{ \langle u, v \rangle - f(u) \right\}, \quad \forall v \in \mathcal{V}$$

The biconjugate function is a minorant of the function

Definition

Let $f: \mathcal{U} \rightarrow \overline{\mathbb{R}}$ be a function

Its **biconjugate** $f^{**'}: \mathcal{U} \rightarrow \overline{\mathbb{R}}$ is defined by

$$f^{**'}(u) = \sup_{v \in \mathcal{V}} \left\{ \langle u, v \rangle - f^*(v) \right\}$$

The inequality below is instrumental in obtaining weak duality

Proposition

For any function $f: \mathcal{U} \rightarrow \overline{\mathbb{R}}$, we have that

$$f^{**'} \leq f$$

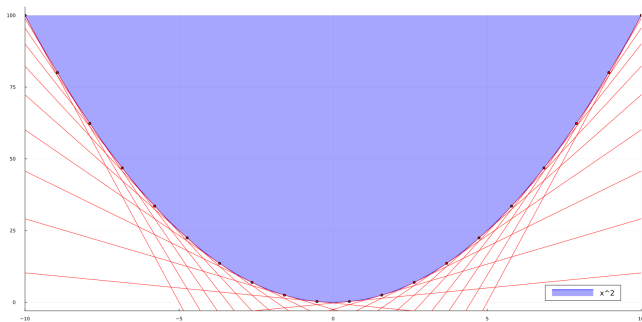
Fenchel-Moreau Theorem

The equality below is instrumental in obtaining strong duality

Theorem

[Fenchel-Moreau]

The function $f: \mathcal{U} \rightarrow \overline{\mathbb{R}}$ is **closed convex** if and only if $f^{**'} = f$



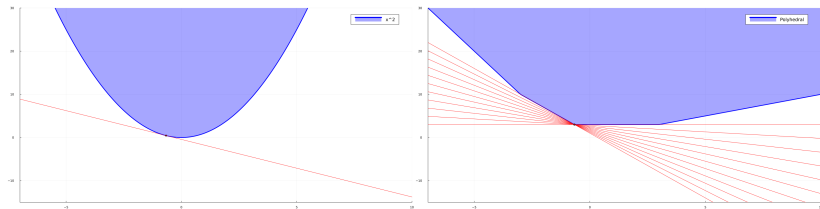
Moreau-Rockafellar subdifferential

Definition

Let $f : \mathcal{U} \rightarrow \overline{\mathbb{R}}$ be a function

Its **subdifferential** $\partial f(u) \subset \mathcal{V}$ at any $u \in \mathcal{U}$ such that $f(u) \in \mathbb{R}$, is defined by

$$v \in \partial f(u) \iff \langle u', v \rangle - f(u') \leq \langle u, v \rangle - f(u), \quad \forall u' \in \mathcal{U}$$



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[Rockafellar, 1974]

1. **Perturb** a minimization problem with a perturbation (primal) variable belonging to a vector space, and a Rockafellian function
2. **Pair** the (primal) perturbation space with a dual space by means of a bilinear form \langle , \rangle
3. **Biconjugate** the perturbation function, and get
 - ▶ a dual problem
 - ▶ weak duality
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Perturbation duality scheme [Rockafellar, 1974]

sets	optimization set \mathcal{X}	primal space \mathcal{U}	pairing $\mathcal{U} \overset{\langle \cdot, \cdot \rangle}{\leftrightarrow} \mathcal{V}$	dual space \mathcal{V}
variables	decision $x \in \mathcal{X}$	perturbation $u \in \mathcal{U}$	$\langle u, v \rangle$ $\in \mathbb{R}$	sensitivity $v \in \mathcal{V}$
bivariate functions		Rockafellian $\mathcal{R}: \mathcal{X} \times \mathcal{U} \rightarrow \overline{\mathbb{R}}$		Lagrangian $\mathcal{L}: \mathcal{X} \times \mathcal{V} \rightarrow \overline{\mathbb{R}}$
definition				$\mathcal{L}(x, v) =$ $\inf_{u \in \mathcal{U}} \{ \mathcal{R}(x, u) - \langle u, v \rangle \}$
property				$-\mathcal{L}(x, \cdot) = (\mathcal{R}(x, \cdot))^*$
property				$-\mathcal{L}(x, \cdot)$ is \star' -convex (hence $\mathcal{L}(x, \cdot)$ is concave usc)
univariate functions		perturbation function $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$		dual function $\psi: \mathcal{V} \rightarrow \overline{\mathbb{R}}$
definition		$\varphi(u) = \inf_{x \in \mathcal{X}} \mathcal{R}(x, u)$		$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v)$
property				$-\psi = \varphi^*$

Weak duality

Perturbation/Rockafellian (Step 1)

Data: set \mathcal{X} , function $f: \mathcal{X} \rightarrow \overline{\mathbb{R}}$ and

original minimization problem $\inf_{x \in \mathcal{X}} f(x)$

- ▶ Embedding/**perturbation scheme** given by a **vector space** \mathcal{U} , and a **Rockafellian** $\mathcal{R}: \mathcal{X} \times \mathcal{U} \rightarrow \overline{\mathbb{R}}$ such that

$$f(x) = \mathcal{R}(x, 0), \quad \forall x \in \mathcal{X}$$

- ▶ The **perturbation function** $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$ is defined by

$$\varphi(u) = \inf_{x \in \mathcal{X}} \mathcal{R}(x, u)$$

original minimization problem $\varphi(0) = \inf_{x \in \mathcal{X}} f(x)$

Duality/Fenchel conjugacy (Steps 2,3)

- Dual vector space \mathcal{V} paired to \mathcal{U} by a bilinear form $\langle \cdot, \cdot \rangle$

We obtain weak duality

$$\begin{aligned}\varphi^{**'}(0) &= \overbrace{\sup_{v \in \mathcal{V}} \{-\varphi^*(v)\}}^{\text{dual problem}} \\ &\leq \\ \varphi(0) &= \underbrace{\inf_{x \in \mathcal{X}} f(x)}_{\text{original problem}}\end{aligned}$$

Weak duality and Lagrangian

Lagrangian

- ▶ Lagrangian $\mathcal{L}: \mathcal{X} \times \mathcal{V} \rightarrow \overline{\mathbb{R}}$ defined by

$$\mathcal{L}(x, v) = \inf_{u \in \mathcal{U}} \left\{ \underbrace{\mathcal{R}(x, u)}_{\text{Rockafellian}} - \langle u, v \rangle \right\}, \quad \forall (x, v) \in \mathcal{X} \times \mathcal{V}$$

- ▶ As $\mathcal{L}(x, v) \leq \mathcal{R}(x, 0) - \langle 0, v \rangle = f(x)$, we get that

$$\sup_{v \in \mathcal{V}} \mathcal{L}(x, v) \leq f(x)$$

hence that

original minimization problem

$$\inf_{x \in \mathcal{X}} \sup_{v \in \mathcal{V}} \mathcal{L}(x, v) \leq \inf_{x \in \mathcal{X}} f(x)$$

Dual function

- ▶ The **dual function** $\psi: \mathcal{V} \rightarrow \overline{\mathbb{R}}$ is defined by

$$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v), \quad \forall v \in \mathcal{V}$$

- ▶ and the **dual problem** is

$$\varphi^{**'}(0) = \sup_{v \in \mathcal{V}} \{ \langle 0, v \rangle - \varphi^*(v) \} = \overbrace{\sup_{v \in \mathcal{V}} \psi(v)}^{\text{dual problem}}$$

$$\begin{aligned} \text{as } -\varphi^*(v) &= -\left(\inf_{x \in \mathcal{X}} \mathcal{R}(x, \cdot) \right)^*(v) \\ &= -\sup_{x \in \mathcal{X}} \left\{ \sup_{u \in \mathcal{U}} \{ \langle u, v \rangle - \mathcal{R}(x, u) \} \right\} \\ &= -\sup_{x \in \mathcal{X}} \left\{ -\underbrace{\inf_{u \in \mathcal{U}} \{ -\langle u, v \rangle + \mathcal{R}(x, u) \}}_{\text{Lagrangian}} \right\} \\ &= \inf_{x \in \mathcal{X}} \mathcal{L}(x, v) = \psi(v) \end{aligned}$$

Weak duality with Lagrangian

$$\begin{aligned} & \varphi^{\star\star'}(0) \\ &= \sup_{v \in \mathcal{V}} \{ -\varphi^\star(v) \} \\ &= \underbrace{\sup_{v \in \mathcal{V}} \inf_{x \in \mathcal{X}} \mathcal{L}(x, v)}_{\text{dual problem}} \\ &\leq \\ &\quad \inf_{x \in \mathcal{X}} \sup_{v \in \mathcal{V}} \mathcal{L}(x, v) \\ &\leq \inf_{x \in \mathcal{X}} f(x) \\ &= \varphi(0) \end{aligned}$$

Strong duality

Strong duality

$$\begin{array}{ccc} \text{dual problem} & & \text{original problem} \\ \underbrace{\varphi^{**'}(0)} & \leq & \underbrace{\varphi(0)} \\ \text{weak duality} & & \end{array}$$

Definition

Strong duality $\iff \varphi^{**'}(0) = \varphi(0) \iff \varphi$ is \star -convex at 0

Paths to strong duality in the convex case

- ▶ Suppose that the **Rockafellian** $\mathcal{R}: \mathcal{X} \times \mathcal{U} \rightarrow \overline{\mathbb{R}}$ is a **jointly convex** function
- ▶ Then, the **perturbation function** $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$ is **convex** as the **marginal function** $\varphi(u) = \inf_{x \in \mathcal{X}} \mathcal{R}(x, u)$
- ▶ If, in addition,
 - ▶ either (global property) the function φ is **proper and lower semicontinuous**, and then $\varphi^{**'} = \varphi$ by the Fenchel-Moreau Theorem,
 - ▶ or (local property) the subdifferential $\partial\varphi(0) \neq \emptyset$, and then the function φ is \star -convex at 0,

and we get **strong duality** $\underbrace{\varphi^{*'}(0)}_{\text{dual problem}} = \underbrace{\varphi(0)}_{\text{original problem}}$

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PDS for convex programming

Background on duality in convex analysis

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Examples of PDS for convex programs

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Conclusion

Lagrangian duality (the case of inequality constraints)

Classic Lagrangian duality (the case of inequality constraints)

- ▶ Optimization set \mathcal{X}
- ▶ Objective function $f: \mathcal{X} \rightarrow]-\infty, +\infty]$
- ▶ Mapping $\theta = (\theta_1, \dots, \theta_m): \mathcal{X} \rightarrow \mathbb{R}^m$, and $\bar{u} \in \mathbb{R}^m$

We consider the optimization problem

$$\min_{\theta(x) \leq \bar{u}} f(x) = \min_{\substack{\theta_1(x) \leq \bar{u}_1 \\ \theta_m(x) \leq \bar{u}_m}} f(x)$$

Perturbation and Rockafellian

- ▶ Perturbation space $\mathcal{U} = \mathbb{R}^m$
- ▶ Rockafellian $\mathcal{R}: \mathcal{X} \times \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\mathcal{R}(x, u) = f(x) + \iota_{\{\theta(x) - \bar{u} \leq u\}} = f(x) + \sum_{j=1}^m \iota_{\{\theta_j(x) - \bar{u}_j \leq u_j\}}$$

Duality, Lagrangian and dual function

- ▶ Dual space $\mathcal{V} = \mathbb{R}^m$
- ▶ We deduce the **Lagrangian** $\mathcal{L}: \mathcal{X} \times \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\begin{aligned}\mathcal{L}(x, v) &= \left(f(x) \dot{+} (-\iota_{\mathbb{R}_+^m}(v)) \right) + \langle \theta(x) - \bar{u}, v \rangle \\ &= \left(f(x) \dot{+} (-\iota_{\mathbb{R}_+^m}(v)) \right) + \sum_{j=1}^m v_j (\theta_j(x) - \bar{u})\end{aligned}$$

- ▶ We deduce the **dual function** $\psi: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v) = (-\iota_{\mathbb{R}_+^m}(v)) \dot{+} \inf_{x \in \mathcal{X}} \left\{ f(x) + \sum_{j=1}^m v_j (\theta_j(x) - \bar{u}) \right\}$$

which is concave upper semicontinuous,
as the supremum of affine functions

Paths to strong duality in the convex case

- ▶ Suppose that
 - ▶ the optimization set \mathcal{X} is a vector space
 - ▶ the objective function $f: \mathcal{X} \rightarrow]-\infty, +\infty]$ is convex
 - ▶ each component of the mapping $\theta = (\theta_1, \dots, \theta_m): \mathcal{X} \rightarrow \mathbb{R}^m$ is a convex function
- ▶ Then, the **perturbation function** $\varphi: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$ is a **convex** function as the marginal

$$\varphi(u) = \inf_{x \in \mathcal{X}} \left\{ f(x) + \sum_{j=1}^m v_j (\theta_j(x) - u) \right\}$$

- ▶ If, in addition,
 - ▶ either the function φ is **proper and lower semicontinuous**
 - ▶ or its subdifferential $\partial\varphi(0) \neq \emptyset$
- then we get **strong duality**

Fenchel-Rockafellar duality

The Fenchel-Rockafellar dual problem

Proposition

adapted from [Rockafellar, 1970, Corollary 31.2.1]

Let $f, g: \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ be proper convex functions
and let $L: \mathbb{R}^n \rightarrow \mathbb{R}^m$ be a linear mapping

$$\sup_{v \in \mathbb{R}^m} \{-g^*(v) - f^*(-L^T v)\} \leq \inf_{x \in \mathbb{R}^n} \{f(x) + g(Lx)\}$$

Furthermore, equality is achieved if either

- ▶ $\exists x \in \text{ri}(\text{dom} f)$ s.t. $Lx \in \text{ri}(\text{dom} g)$
- ▶ $\exists v \in \text{ri}(\text{dom} g^*)$ s.t. $L^T v \in \text{ri}(\text{dom} f^*)$

Perturbation and Rockafellian

- ▶ Perturbation space $\mathcal{U} = \mathbb{R}^m$
- ▶ Rockafellian $\mathcal{R}: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\mathcal{R}(x, u) = f(x) + g(Lx + u)$$

Duality and Lagrangian and dual function

- ▶ Dual space $\mathcal{V} = \mathbb{R}^m$
- ▶ We deduce the **Lagrangian** $\mathcal{L}: \mathcal{X} \times \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\mathcal{L}(x, v) = \langle x, L^T v \rangle + f(x) - g^*(v)$$

- ▶ We deduce the **dual function** $\psi: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v) = -g^*(v) + \underbrace{\left(- \sup_{x \in \mathcal{X}} \{ \langle x, -L^T v \rangle - f(x) \} \right)}_{f^*(-L^T v)}$$

Application to regularized problems

For a linear mapping $L: \mathbb{R}^n \rightarrow \mathbb{R}^m$ and a proper convex function $f: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$ and suitable assumptions

$$\sup_{v \in \mathbb{R}^m} -\|L^T v\|^2 - f^*(v) = \inf_{x \in \mathbb{R}^n} f(Lx) + \frac{1}{2} \|x\|^2$$

Can be useful for computation if $m < n$ and f^* easy to compute

Semidefinite programming dual problem

adapted from [Calafiore and El Ghaoui, 2014, Chapter 11]

Let \mathbb{S}^n be the set of $n \times n$ symmetric matrices

Let $\mathbb{S}_+^n \subset \mathbb{S}^n$ be the set of $n \times n$ semidefinite matrices

Proposition

Let $K, A_1, \dots, A_m \in \mathbb{S}^n$ and $b \in \mathbb{R}^m$

Then, we have

$$\sup_{\substack{v \in \mathbb{R}^m \\ K - \sum_{j=1}^m v_j A_j \succeq 0}} \langle b, v \rangle \leq \inf_{\substack{X \in \mathbb{S}^n \\ \text{trace}(A_j X) = b_j, \ j=1, \dots, m \\ X \succeq 0}} \text{trace}(KX)$$

Furthermore, equality is achieved
if some Slater's condition is satisfied

Perturbation and Rockafellian

- ▶ Perturbation space $\mathcal{U} = \mathbb{R}^m$
- ▶ Rockafellian $\mathcal{R}: \mathbb{S}^n \times \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\mathcal{R}(x, u) = \text{trace}(KX) + \iota_{X \succeq 0} + \sum_{j=1}^m \iota_{\text{trace}(A_j X) = b_j + u_j}$$

where $u = \begin{pmatrix} u_1 \\ \vdots \\ u_m \end{pmatrix}$

Duality and Lagrangian and dual function

- ▶ Dual space $\mathcal{V} = \mathbb{R}^m$
- ▶ We deduce the **Lagrangian** $\mathcal{L}: \mathcal{X} \times \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\mathcal{L}(X, v) = \langle b, v \rangle + \text{trace}\left((K - \sum_{j=1}^m v_j A_j)X\right)$$

- ▶ We deduce the **dual function** $\psi: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

$$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v) = \langle b, v \rangle - \iota_{K - \sum_{j=1}^m v_j A_j \succeq 0}$$

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Introducing generalized convexity

Fenchel conjugate $f^*(v) = \sup_{u \in \mathbb{R}^m} \langle u, v \rangle - f(u)$	c-conjugate $g^c(v) = \sup_{u \in U} c(u, v) \dot{+} (-g(u))$
Fenchel biconjugate $f^{**'}(u) = \sup_{v \in \mathbb{R}^m} \langle u, v \rangle - f^*(v)$	c-biconjugate $g^{cc'}(u) = \sup_{v \in V} c(u, v) \dot{+} (-g^c(v))$
\star - convex functions $\iff f^{**'} = f$	c-convex functions $\iff g^{cc'} = g$

with the Moreau lower and upper additions

$$(+\infty) \dot{+} (-\infty) = (-\infty) \dot{+} (+\infty) = -\infty$$

$$(+\infty) \dot{+} (-\infty) = (-\infty) \dot{+} (+\infty) = +\infty$$

Generalized perturbation-duality scheme [Balder, 1977]

sets	optimization set \mathcal{X}	primal set \mathcal{U}	coupling $\mathcal{U} \leftrightarrow \mathcal{V}$	dual set \mathcal{V}
variables	decision $x \in \mathcal{X}$	perturbation $u \in \mathcal{U}$	$c(u, v)$ $\in \mathbb{R}$	sensitivity $v \in \mathcal{V}$
bivariate functions		Rockafellian $\mathcal{R}: \mathcal{X} \times \mathcal{U} \rightarrow \overline{\mathbb{R}}$		Lagrangian $\mathcal{L}: \mathcal{X} \times \mathcal{V} \rightarrow \overline{\mathbb{R}}$
definition				$\mathcal{L}(x, v) =$ $\inf_{u \in \mathcal{U}} \{ \mathcal{R}(x, u) + (-c(u, v)) \}$
property				$-\mathcal{L}(x, \cdot) = (\mathcal{R}(x, \cdot))^c$
property				$-\mathcal{L}(x, \cdot)$ is c' -convex
univariate functions		perturbation function $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$		dual function $\psi: \mathcal{V} \rightarrow \overline{\mathbb{R}}$
definition		$\varphi(u) = \inf_{x \in \mathcal{X}} \mathcal{R}(x, u)$		$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v)$
property				$-\psi = \varphi^c$

- ▶ Anchor $\bar{u} \in \mathcal{U}$ and dual maximization problem (weak duality)
 $\varphi^{cc'}(\bar{u}) = \sup_{v \in \mathcal{V}} \{ c(\bar{u}, v) + \psi(v) \} \leq \inf_{x \in \mathcal{X}} f(x) = \varphi(\bar{u})$
- ▶ Strong duality iff φ is c -convex at \bar{u} iff $\varphi^{cc'}(\bar{u}) = \varphi(\bar{u})$

Case of evaluation couplings

Case of evaluation couplings (developped later)

- ▶ Given a primal set \mathcal{U} and a function set $\mathcal{F} \subset \{F: \mathcal{U} \rightarrow \overline{\mathbb{R}}\}$, the **evaluation coupling** $c_{\mathcal{F}}: \mathcal{U} \times \mathcal{F} \rightarrow \overline{\mathbb{R}}$ is defined by

$$c_{\mathcal{F}}(u, F) = F(u), \quad \forall u \in \mathcal{U}, F \in \mathcal{F}$$

- ▶ For a given (perturbation) function $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$, **weak duality** is always achieved

$$\varphi^{c_{\mathcal{F}}c_{\mathcal{F}'}} \leq \varphi$$

- ▶ Sufficient condition for **strong duality**

$$\varphi \in \mathcal{F} \implies \varphi^{c_{\mathcal{F}}c_{\mathcal{F}'}} = \varphi$$

- ▶ Two **trivial cases** of strong duality

1. $\mathcal{F} = \{F: \mathcal{U} \rightarrow \overline{\mathbb{R}}\} = \overline{\mathbb{R}}^{\mathcal{U}}$
2. $\mathcal{F} = \{\varphi\}$

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Pure integer linear program in standard form

$$\begin{aligned} \inf \quad & k^T x \\ x \in & \mathbb{Z}^n \\ Ax = & b \\ x \geq & 0 \end{aligned}$$

Usual continuous dual in PILP

$$\begin{array}{ccc} \text{dual (continuous) problem} & & \text{primal (integer) problem} \\ \hline \sup_{\substack{\lambda \in \mathbb{R}^m \\ A^T \lambda \leq k}} \langle b \mid \lambda \rangle & \leq & \inf_{\substack{x \in \mathbb{Z}^n \\ Ax = b \\ x \geq 0}} k^T x \\ & \text{weak duality} & \end{array}$$

- ▶ Right-hand side b perturbation and scalar product coupling
- ▶ Usually strong duality is not achieved
- ▶ Can we design tighter dual problems?
(Useful for Branch-and-bound like methods)

Changing the perturbation/changing the coupling

$$\varphi(b) = \inf_{\substack{x \in \mathbb{Z}_+^n \\ Ax = b}} k^T x$$

$$\langle, \rangle : \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$$

$$\varphi(b) = \dots$$

$$\langle, \rangle : \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$$

$$\varphi(b) = \inf_{\substack{x \in \mathbb{Z}_+^n \\ Ax = b}} k^T x$$

$$c : \mathbb{R}^m \times \mathcal{F} \rightarrow \overline{\mathbb{R}}$$

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(Geoffrion) Lagrangian relaxation [Geoffrion, 1974]

1. We partially perturb

$$\forall b^1 \in \mathbb{R}^{m_1}, \quad \varphi(b^1) = \inf_x \quad k^T x$$
$$A^1 x = b^1$$
$$A^2 x = b^2$$
$$x \geq 0$$
$$x \in \mathbb{Z}^n$$

2. We pair the primal space \mathbb{R}^{m_1} and a dual space \mathbb{R}^{m_1}

$$\langle, \rangle : \mathbb{R}^{m_1} \times \mathbb{R}^{m_1} \rightarrow \mathbb{R}$$

3. We biconjugate the perturbation function φ

$$\varphi^{**'}(b^1) = \sup_{\lambda \in \mathbb{R}^{m_1}} \underbrace{\inf_{\substack{A^2 x = b^2 \\ x \geq 0 \\ x \in \mathbb{Z}^n}} k^T x + \langle b^1 - A^1 x \mid \lambda \rangle}_{g(\lambda)}$$

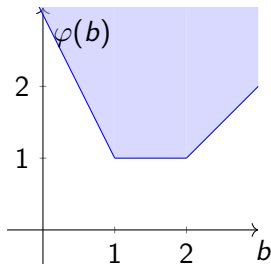
Example of perturbation function's epigraph for LP

Example

Let $\varphi : \mathbb{R} \rightarrow \mathbb{R}$ be defined as

$$\begin{aligned} \varphi(b) = \quad & \inf_{x \in \mathbb{R}^3} && x_1 + x_2 + x_3 \\ & x_1 + x_2 + 3x_3 = 1 \\ & x_1 + 2x_2 + 4x_3 = b \\ & x \geq 0 \end{aligned}$$

Then $\varphi(b) = \max\{3 - 2b, 1, b - 1\}$, $\forall b \in \mathbb{R}$



Condition for tighter gap than continuous dual problem

adapted from [Conforti, Cornuéjols, and Zambelli, 2014, Corollary 8.4.]

Proposition

If A^2 and b^2 are **rational** then

$$\sup_{\substack{\lambda \in \mathbb{R}^m \\ A^T \lambda \leq k}} \langle b \mid \lambda \rangle \leq \varphi^{**'}(b^1) = \sup_{\lambda \in \mathbb{R}^{m_1}} g(\lambda)$$

where

$$g(\lambda) = \inf_{\substack{A^2 x = b^2 \\ x \geq 0 \\ x \in \mathbb{Z}^n}} k^T x + \langle b^1 - A^1 x \mid \lambda \rangle$$

$$\text{and } A = \begin{pmatrix} A^1 \\ A^2 \end{pmatrix}, \quad b = \begin{pmatrix} b^1 \\ b^2 \end{pmatrix}, \quad m = m_1 + m_2$$

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Evaluation coupling

Definition

Let $\mathcal{F} \subset \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$ be a set of functions

We call $c_{\mathcal{F}}: \mathbb{R}^m \times \mathcal{F} \rightarrow \overline{\mathbb{R}}$ defined by

$$c_{\mathcal{F}}(b, F) = F(b), \quad \forall b \in \mathbb{R}^m, \forall F \in \mathcal{F}$$

the **evaluation coupling** of \mathcal{F}

Remark

- ▶ Here the **dual variables are functions**
- ▶ If $\mathcal{F} = \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}} \mid F \text{ is linear}\}$, then $c_{\mathcal{F}} = \langle | \rangle$

Resulting evaluation dual problem

also see [Tind and Wolsey, 1981, Sect. 6]

Consider the perturbation function defined by

$$\forall b \in \mathbb{R}^m, \quad \varphi(b) = \inf_{\substack{x \\ Ax = b \\ x \geq 0 \\ x \in \mathbb{Z}^n}} k^T x$$

Proposition

Let $\mathcal{F} \subset \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$ be a set of functions

Then, for any $b \in \mathbb{R}^m$

$$\varphi^{c_{\mathcal{F}} c_{\mathcal{F}'}}(b) = \sup_{F \in \mathcal{F}} \left\{ F(b) + \inf_{x \in \mathbb{Z}_+^n} \{k^T x - F(Ax)\} \right\}$$

$$\varphi^{c_{\mathcal{F}} c_{\mathcal{F}'}}(b) \underbrace{\leq}_{\text{weak duality}} \varphi(b)$$

weak duality

Proof. Compute first conjugate

$$\begin{aligned}\varphi^c(F) &= \sup_{b \in \mathbb{R}^m} \{c(b, F) + (-\varphi(b))\} \\&= \sup_{b \in \mathbb{R}^m} \left\{ c(b, F) + \left(- \inf_{\substack{x \in \mathbb{Z}_+^n \\ Ax=b}} k^T x \right) \right\} \\&= \sup_{b \in \mathbb{R}^m} \left\{ c(b, F) + \sup_{\substack{x \in \mathbb{Z}_+^n \\ Ax=b}} -k^T x \right\} \\&= \sup_{x \in \mathbb{Z}_+^n} \left\{ -k^T x + \sup_{\substack{b \in \mathbb{R}^m \\ Ax=b}} c(b, F) \right\} \\&= \sup_{x \in \mathbb{Z}_+^n} \left\{ -k^T x + c(Ax, F) \right\} \\&= - \inf_{x \in \mathbb{Z}_+^n} \left\{ \underbrace{k^T x - F(Ax)}_{\text{Lagrangian } \mathcal{L}(x, F)} \right\}\end{aligned}$$

Revisiting the Fenchel coupling with evaluation coupling

- ▶ Bilinear coupling $\langle \cdot | \cdot \rangle : \overbrace{\mathbb{R}^m}^{\text{perturbation space}} \times \overbrace{\mathbb{R}^m}^{\text{dual space}} \rightarrow \mathbb{R}$
- ▶ \mathbb{R}^m can be identified to the functional space $\Lambda = \{F: \mathbb{R}^m \rightarrow \mathbb{R} | F \text{ is linear}\}$

$$\lambda \in \mathbb{R}^m \leftrightarrow F \in \Lambda$$

$$\langle b | \lambda \rangle \leftrightarrow c_\Lambda(b, F)$$

- ▶ Thus, the resulting dual problem

$$\begin{aligned} \varphi^{c_\Lambda c'_\Lambda}(b) &= \sup_{\lambda \in \mathbb{R}^m} \left\{ \langle b | \lambda \rangle + \underbrace{\inf_{x \in \mathbb{Z}_+^n} \{k^T x - \langle Ax | \lambda \rangle\}}_{\iota_{A^T \lambda \leq k}} \right\} \\ &= \sup_{\substack{\lambda \in \mathbb{R}^m \\ A^T \lambda \leq k}} \langle b | \lambda \rangle \end{aligned}$$

Inclusion of functional sets

Consider the perturbation function defined by

$$\forall b \in \mathbb{R}^m, \quad \varphi(b) = \inf_x \quad k^T x$$
$$\begin{aligned} Ax &= b \\ x &\geq 0 \\ x &\in \mathbb{Z}^n \end{aligned}$$

Proposition

If $\mathcal{F}^1 \subset \dots \subset \mathcal{F}^J \subset \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$, then

$$\varphi^{c_{\mathcal{F}^1} c_{\mathcal{F}^1}'} \leq \dots \leq \varphi^{c_{\mathcal{F}^J} c_{\mathcal{F}^J}'} \leq \varphi$$

The **larger** the set of dual functions, the **tighter** the gap!

Characterization of strong duality for evaluation coupling

Proposition

[Tind and Wolsey, 1981, Proposition 6.8]

Let $\mathcal{F} \subset \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$ be a set of functions and $f: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

Then

$$f^{c_{\mathcal{F}} c_{\mathcal{F}'}} = f \iff \exists \{f_i\}_{i \in I} \subset \mathcal{F} \text{ s.t. } f = \sup_{i \in I} f_i$$

Remark

Cases when the *equality* is trivially true

- ▶ \mathcal{F} is *too "general"*: $\mathcal{F} = \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$
- ▶ \mathcal{F} is *too "specific"*: $\mathcal{F} = \{f\}$

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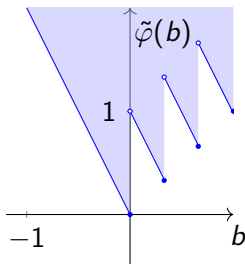
Example

Let $\varphi: \mathbb{R} \rightarrow \mathbb{R}$ be defined as

$$\forall b \in \mathbb{R}, \quad \varphi(b) = \inf_{\substack{x \in \mathbb{Z}^2 \\ x_1 - x_2 = b \\ x \geq 0}} x_1 + 2x_2$$

Then φ coincides on its domain with

$$\tilde{\varphi}(b) = \max\{-2b, \lceil 3b \rceil - 2b\}, \quad \forall b \in \mathbb{R}$$



Summary of the perturbation-duality scheme for PILP

1. We **perturb** a minimization problem

$$\forall \mathbf{b} \in \mathbb{R}^m, \quad \varphi(\mathbf{b}) = \inf_{\substack{\mathbf{x} \\ A\mathbf{x} = \mathbf{b} \\ \mathbf{x} \in \mathbb{Z}_+^n}} k^T \mathbf{x}$$

2. We pair the primal space \mathbb{R}^m and a function set \mathcal{F}

$$\begin{aligned} c_{\mathcal{F}}: \mathbb{R}^m \times \mathcal{F} &\rightarrow \overline{\mathbb{R}} \\ c_{\mathcal{F}}(\mathbf{b}, F) &= F(\mathbf{b}) \end{aligned}$$

3. We biconjugate the perturbation function φ

$$\begin{aligned} &\overbrace{\varphi^{cc'}(\mathbf{b}) \leq \varphi(\mathbf{b}), \quad \forall \mathbf{b} \in \mathbb{R}^m}^{\text{weak duality is guaranteed}} \\ \varphi^{cc'}(\mathbf{b}) &= \sup_{F \in \mathcal{F}} \left\{ F(\mathbf{b}) + \inf_{\mathbf{x} \in \mathbb{Z}_+^n} \{ k^T \mathbf{x} - F(A\mathbf{x}) \} \right\} \end{aligned}$$

Strong duality for the subadditive dual problem

Let \mathcal{S} be the set of **subadditive functions**

$$\mathcal{S} = \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}} \text{ s.t. } F(b^1 + b^2) \leq F(b^1) + F(b^2), \forall b^1, b^2\}$$

$$(-\infty) + (+\infty) = (+\infty) + (-\infty) = +\infty$$

Proposition

Let $A \in \mathbb{R}^{m \times n}$ and $k \in \mathbb{R}^n$

Let $c_{\mathcal{S}}: \mathbb{R}^m \times \mathcal{S} \rightarrow \overline{\mathbb{R}}$ be the evaluation coupling of \mathcal{S}

Then, for all $b \in \mathbb{R}^m$

$$\varphi^{c_{\mathcal{S}} c_{\mathcal{S}}'}(b) = \varphi(b)$$

Proof: φ is subadditive

Proof of subadditive strong duality. φ is subadditive

(i) Let $b^1, b^2 \in \mathbb{R}^m$

(ii) ▶ Assume there are $x^1, x^2 \in \mathbb{Z}_+^n$ s.t.

$$Ax^1 = b^1, \quad Ax^2 = b^2$$

Then $A(x^1 + x^2) = b^1 + b^2$

▶ For all such x^1, x^2 and by definition of φ

$$\varphi(b^1 + b^2) \leq \langle k \mid x^1 + x^2 \rangle = \langle k \mid x^1 \rangle + \langle k \mid x^2 \rangle$$

▶ Going to the infimum in x^1 , then in x^2

$$\varphi(b^1 + b^2) \leq \varphi(b^1) + \varphi(b^2)$$

(iii) If there is no such $x^1, x^2 \in \mathbb{Z}_+^n$

Then $\varphi(b^1) = +\infty$ or $\varphi(b^2) = +\infty$

Thus, by definition of $\dot{+}$

$$\varphi(b^1 + b^2) \leq \varphi(b^1) \dot{+} \varphi(b^2)$$

Previous work on PILP duality

- ▶ **Surveys** of previous works
[Tind and Wolsey, 1981], [Güzelsoy, Ralphs, and Cochran, 2010]
- ▶ Some pioneer papers on strong duality in **the rational case**
(when $A \in \mathbb{Q}^{m \times n}$, $k \in \mathbb{Q}^n$)
[Johnson, 1973], [Jeroslow, 1979], [Wolsey, 1981], [Blair and Jeroslow, 1982]
- ▶ **Computation** of optimal dual functions
[Wolsey, 1981], [Klabjan, 2007]

Subadditive dual problem in [Jeroslow, 1979]

$$\begin{aligned} & \sup && F(\bar{b}) \\ & F: \mathbb{R}^m \rightarrow \mathbb{R} \\ & F(A_j) \leq k_j \\ & F(0) \leq 0 \\ & F \text{ is subadditive} \end{aligned}$$

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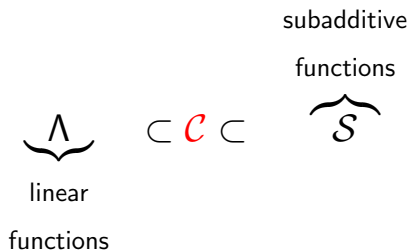
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Somewhere in between the linear and the subadditive



Definition of Chvátal functions

Definition

The class of Chvátal functions \mathcal{C} is the smallest class of functions $\mathcal{F} \subset \{F: \mathbb{R}^m \rightarrow \mathbb{R}\}$ such that

$$b \in \mathbb{R}^m \mapsto \langle b \mid \lambda \rangle \in \mathcal{F}, \quad \forall \lambda \in \mathbb{Q}^m \quad (\text{linear functions})$$

$$\alpha F_1 + \beta F_2 \in \mathcal{F}, \quad \forall F_1, F_2 \in \mathcal{F}, \quad \alpha, \beta \in \mathbb{Q}_+ \quad (\text{conic combination})$$

$$\lceil F \rceil \in \mathcal{F}, \quad \forall F \in \mathcal{F} \quad (\text{round-up})$$

Examples in 1D

- ▶ $b \mapsto \frac{3}{4}b$
- ▶ $b \mapsto \lceil b \rceil$
- ▶ $b \mapsto \frac{3}{4}b + \frac{7}{10}\lceil b \rceil$
- ▶ $b \mapsto 15b + \frac{39}{22}\left[\frac{3}{4}b + \frac{7}{10}\lceil b \rceil\right] + \lceil 16b \rceil$

Strong duality with Chvátal functions

adapted from [Blair and Jeroslow, 1982]

We define a perturbation function

$$\forall b \in \mathbb{R}^m, \quad \varphi(b) = \inf_{\substack{x \\ Ax = b \\ x \in \mathbb{Z}_+^n}} k^T x$$

Proposition

We remind that c_C is the evaluation coupling of the Chvátal functions
If $A \in \mathbb{Q}^{m \times n}$ and $k \in \mathbb{Q}^n$ then

$$\varphi^{c_C c'_C}(b) = \varphi(b), \quad \forall b \in \text{dom} \varphi$$

Remark

The perturbation function φ is defined on \mathbb{R}^m but $\text{dom} \varphi \subset \mathbb{Q}^m$

Outline of the presentation

PDS for linear programming (LP)

PDS for convex programming

PDS for pure integer linear programming (PILP)

Conclusion

Steps of the perturbation-duality scheme

[Rockafellar, 1974]

1. **Perturb** a minimization problem with a perturbation (primal) variable belonging to a vector space, and a Rockafellian function
2. **Pair** the (primal) perturbation space with a dual space by means of a bilinear form $\langle | \rangle$
3. **Biconjugate** the perturbation function, and get
 - ▶ a dual problem
 - ▶ weak duality
4. Deduce conditions for **strong duality** by means of either global or local properties of the perturbation function

Branching out: rank restricted Chvátal functions

Rank restricted Chvátal functions

Let $F \in \mathcal{C}$ be a Chvátal function

Definition

The **rank** of F is defined *informally* as the smallest number of $\lceil \cdot \rceil$ needed to encode F

We denote by $\mathcal{C}^r \subset \mathcal{C}$ the Chvátal function of rank **not greater** than $r \in \mathbb{N}$

- Inclusion of function sets

$$\underbrace{\mathcal{C}^0}_{\text{linear functions with rational } \lambda} \subset \mathcal{C}^1 \subset \dots \subset \mathcal{C}^m \subset \dots \subset \mathcal{S}$$

linear functions
with rational λ

- Weak duality chain

$$\varphi^{\mathcal{C}^0 \mathcal{C}^0'} \leq \varphi^{\mathcal{C}^1 \mathcal{C}^1'} \leq \dots \varphi^{\mathcal{C}^r \mathcal{C}^r'} \leq \dots \leq \varphi^{\mathcal{C}^S \mathcal{C}^S'} \leq \varphi$$

Example of partially perturbed restricted Chvátal scheme

- ▶ We define a perturbation function

$$\forall b \in \mathbb{R}^{m_1}, \varphi(b^1) = \inf_x k^T x$$
$$\begin{aligned} A^1 x &= b^1 \\ A^2 x &= b^2 \\ x &\geq 0 \\ x &\in \mathbb{Z}^n \end{aligned}$$

- ▶ We define a coupling between primal and dual space

$$c : \mathbb{R}^{m_1} \times \mathbb{Q}^{m_1} \times \mathbb{Q}_+ \rightarrow \mathbb{R} \text{ (for given } \beta \in \mathbb{Q}^{m_1})$$
$$c(b^1, (\lambda, \alpha)) = \lambda^T b^1 + \alpha \left[\beta^T b^1 \right],$$
$$\forall b^1 \in \mathbb{R}^{m_1}, \forall (\lambda, \alpha) \in \mathbb{Q}^{m_1} \times \mathbb{Q}_+$$

- ▶ We biconjugate the perturbation functions

$$\underbrace{\varphi^{cc'}(b^1) \leq \varphi(b^1), \forall b \in \mathbb{R}^{m_1}}_{\text{weak duality}}$$

Resulting dual problems yields a tighter gap

$$\varphi^{cc'}(b) = \sup_{(\lambda, \alpha) \in \mathbb{Q}^m \times \mathbb{Q}_+} \underbrace{\left\{ \lambda^T b^1 + \alpha \left\lceil \beta^T b^1 \right\rceil + \inf_{\substack{A^2 x = b^2 \\ x \geq 0 \\ x \in \mathbb{Z}^n}} \left\{ k^T x - \lambda^T A^1 x + \alpha \left\lceil \beta^T A^1 x \right\rceil \right\} \right\}}_{\tilde{g}(\lambda, \alpha)}$$

We have a tighter gap

$$\underbrace{\sup_{\lambda \in \mathbb{Q}^{m_1}} g(\lambda)}_{\text{Lagrangian relaxation}} \leq \sup_{(\lambda, \alpha) \in \mathbb{Q}^{m_1} \times \mathbb{Q}_+} \tilde{g}(\lambda, \alpha) \leq \underbrace{\varphi(b)}_{\text{original problem}}$$

Thank you for your attention !

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